

# Data Structures for Dynamic and Kinetic Multidimensional Point Sets

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# Motivation

## Latent-Space Embedding for Large Networks

Performed through iterative search. Large networks involve:

- solution spaces of **very high dimension**
- a **large number of iterations** to achieve convergence
- inner-loop computations of **quadratic** complexity

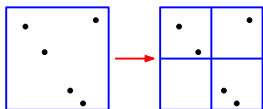
## Recent results in spatial data structures

- Simple and practical **dynamic structures**
- Intrinsic data structures for **kinetic updates**
- Fast approximations through **hierarchical sketching**
- Space/query-time **tradeoffs**

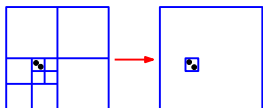
# Basic Structures: Compressed Quadtree

**Quadtree:** Hierarchical structure based on subdivision of squares.

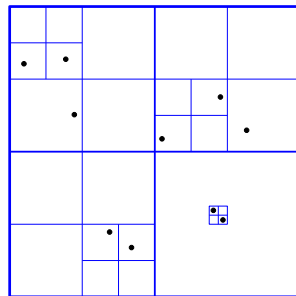
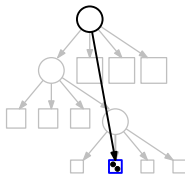
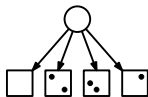
**Path Compression:** Chains of trivial splits are compressed.



splitting



compression (shrinking)



# Basic Structures: Balanced Partition Trees

**Problem:** Compressed quadtrees may have depth  $\Omega(n)$ .

## Balanced Structures

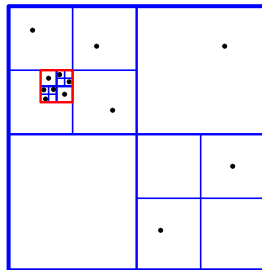
Achieving  $O(\log n)$  depth:

**BBD-tree:** [AMN98] Combines **splitting** with **centroid shrinking**.

Produces an inner cell and outer cell.

**BAR-tree:** [DGK01]

Uses **slanted splitting planes** to split clusters. (Produces convex cells.)



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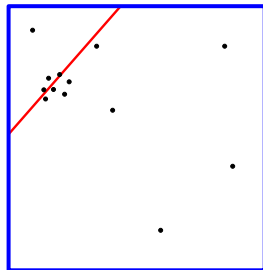
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# Approximate Spherical Range Searching

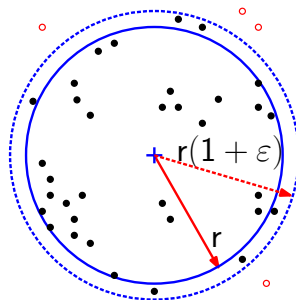
**Example:** Count (or report) the points lying within a given spherical range. Allow an error of  $\epsilon$ .

Preprocessing: Build BBD tree.

Query Processing:

- Find maximal cells lying within the outer range and covering the inner range
- Access counts for each cell
- Return the total

Query time:  $O(\log n + (1/\epsilon)^{d-1})$ .



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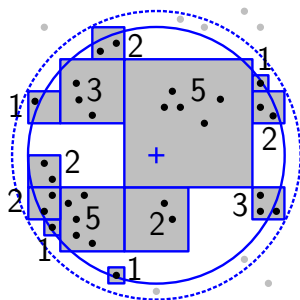
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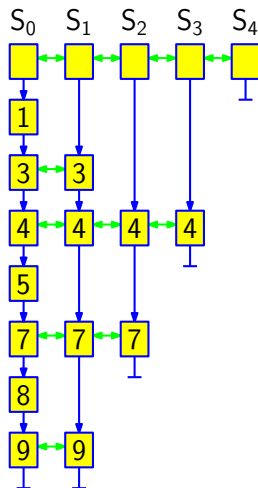
# Dynamic (Balanced) Structures

How to support **insertion** and **deletion**?

## Analogy — Skiplist [Pug90]

A randomized structure for 1-dimensional data. Build linked lists for successive **random samples**.

- $S_0 \leftarrow S$  (the original point set).
- $S_1 \leftarrow$  sample  $S_0$  with probability  $\frac{1}{2}$ .
- $S_2 \leftarrow$  sample  $S_1$  with probability  $\frac{1}{2}$ .
- The process ends after  $O(\log n)$  stages in expectation.





# Skip Quadtree

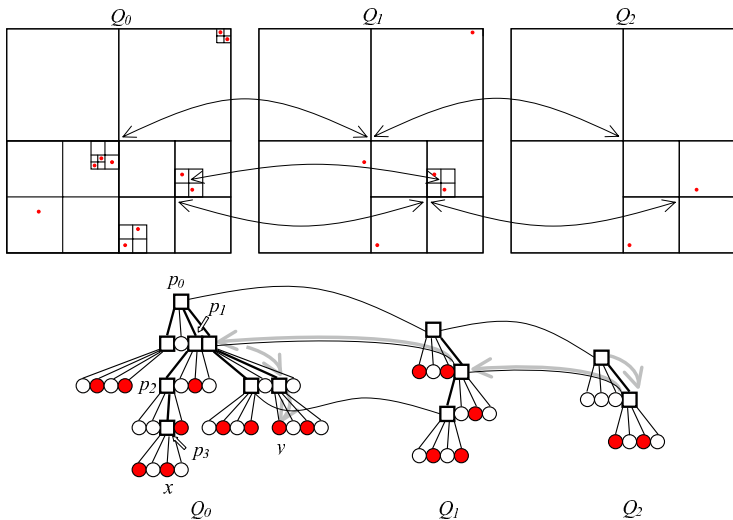
## Skip Quadtree [EGS05]

Same idea, but applied to a sequence of **compressed quadtrees**.

- $Q_0 \leftarrow$  quadtree for  $S_0$ .
- $Q_1 \leftarrow$  quadtree for  $S_1$ .
- $Q_2 \leftarrow$  quadtree for  $S_2$ .
- ... Each node of  $Q_i$  is **linked** to its counterpart in  $Q_{i-1}$ .

Although each quadtree may be unbalanced (like a linked list) it is possible to access each node in  $O(\log n)$  time through the links.

# Skip Quadtree



# Kinetic Structures

Latent embedding algorithms involve adjusting the location many points with each iteration. The motion of each point may be small.

## Kinetic Updates

Update a data structure after a **small motion** involving **many points** of the set.

- Data structures that are defined in terms of a fixed **coordinate frame** are sensitive to changes in **absolute position** of points, even if the **relative position** remains **unchanged**.
- Can we define spatial data structures that are **independent** of any coordinate frame?

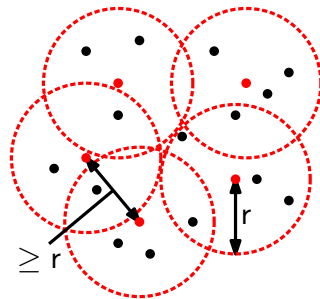
# $r$ -Net

## $r$ -Net

... is a subset  $X \subseteq S$  such that

- (i) every point of  $S$  is within distance  $r$  of some point of  $X$
- (ii) the pairwise distance between any two points of  $X$  is  $\geq r$

The **balls** of radius  $r$  form a cover, which is similar to the partition provided by the **square cells** of one level of a quadtree.

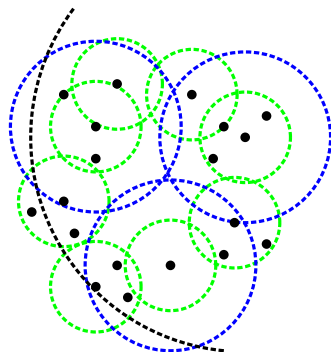


# Net Tree

## Net Tree [CG06, GGN04, HM06]

- The leaves are  $S_0 = S$ .
- Let  $r$  be the minimum distance between any two points of  $S$ .
- Let  $S_1$  be an  $r$ -net of  $S_0$ .
- Let  $S_2$  be a  $(2r)$ -net of  $S_1$ .
- Let  $S_j$  be a  $(2^j r)$ -net of  $S_{j-1}$ .
- ...
- Until only one remains — the root.

Similar to the quadtree in spirit, but intrinsic to the point set.



# Challenges

**Depth vs. Breadth:** Most innovations have focused on maintaining structures of low **depth**. However, search times are dominated by the search's **breadth**, that is, the number of cells visited.

**Tight vs. Slack:** The best static structures achieve efficiency by enforcing **tight constraints** on subdivision properties (e.g. tree depth, cell size, aspect ratio). However, tight constraints result in frequent certificate failures, as used by kinetic structures.

**The Right Mixture:** What is the proper mixture of methods to achieve best overall performance, in terms of accuracy and execution times?

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