Approximate Sampling for Binary Discrete Exponential Families, with Fixed Execution Time and Quality Guarantees

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Exponential Families on Binary Vectors

- Let $Y = Y_1, \ldots, Y_N$ be a finite vector of binary-valued random variables with joint support \mathcal{Y}_N
- \blacktriangleright Wlg, we may write the joint pmf of Y in discrete exponential family form as

$$\Pr\left(Y = y \left|\theta\right.\right) = \frac{\exp\left[\theta^T t\left(y\right)\right]}{\sum_{y' \in \mathcal{Y}_N} \exp\left[\theta^T t\left(y'\right)\right]} I_{\mathcal{Y}_N}(y) \tag{1}$$

- $\triangleright t : \mathcal{Y}_N \mapsto \mathbb{R}^p$ is a vector of sufficient statistics, $\theta \in \mathbb{R}^p$ is a parameter vector, and $I_{\mathcal{Y}_N}$ is an indicator function for membership in the support
- ▷ In general, cannot compute due to intractability of $\sum_{y' \in \mathcal{Y}_N} \exp \left[\theta^T t(y')\right]$
- ► Useful feature of above full conditionals easily obtained:

$$\Pr\left(Y_{i}=1 \mid \theta, Y_{i}^{c}=y_{i}^{c}\right) = \left[1 + \exp\left[\theta^{T}\left[t\left(y_{i}^{-}\right) - t\left(y_{i}^{+}\right)\right]\right]\right]^{-1}$$
(2)

▷ Y_i^c refers to all elements of Y other than the *i*th, Y_i^+ refers to the vector Y w/*i*th entry is set to 1, and Y_i^- refers to Y w/*i*th entry set to 0



• Note that we may also write Y's joint distribution in sequential form:

$$\Pr(Y = y | \theta) = \Pr(Y_1 = y_1 | \theta) \Pr(Y_2 = y_2 | \theta, Y_1 = y_1) \dots \times \Pr(Y_N = y_N | \theta, Y_1 = y_1, \dots, Y_{N-1} = y_{N-1})$$
(3)

But what are these "partial conditionals" in the sequence?

• Let $Y_{\leq i}^c = Y_1, \ldots, Y_{i-1}$. By definition,

$$\Pr\left(Y_{i} = 1 \mid \theta, Y_{$$

So the "partial conditionals" are convex combinations of the full conditionals for Y_i
 And, as noted, those full conditionals are often easy to work with....

Bounding the Partial Conditionals

- An implication of Eq 4: full conditionals can be used to bound partial conditionals (a la Butts (2010))
 - $\triangleright \min_{y' \in \mathcal{Y}_N : y_{<i}^c = y_{<i}^c} \Pr(Y_i = y_i | Y_i^c = y_i'^c) \le \Pr(Y_i = y_i | Y_{<i} = y_{<i})$
 - $\triangleright \max_{y' \in \mathcal{Y}_N : y_{<i}^c = y_{<i}^c} \Pr(Y_i = y_i | Y_i^c = y_i'^c) \ge \Pr(Y_i = y_i | Y_{<i} = y_{<i})$
 - \triangleright Minima/maxima are over all sequences preserving outcomes prior to Y_i
 - Follows immediately from convexity condition
- ▶ Using Eq 2, we can easily express the bounds in logit form
 - Define the following:

$$\diamond \alpha_{i} \leq \min_{y' \in \mathcal{Y}_{N}: y_{\leq i}^{c} = y_{\leq i}^{c}} \left[1 + \exp\left[\theta^{T}\left[t\left(y_{i}^{\prime-}\right) - t\left(y_{i}^{\prime+}\right)\right]\right] \right]^{-1}$$

$$\diamond \beta_{i} \equiv \Pr(Y_{i} = 1 | Y_{\leq i} = y_{\leq i})$$

$$\diamond \gamma_{i} \geq \max_{y' \in \mathcal{Y}_{N}: y_{\leq i}^{\prime c} = y_{\leq i}^{c}} \left[1 + \exp\left[\theta^{T}\left[t\left(y_{i}^{\prime-}\right) - t\left(y_{i}^{\prime+}\right)\right]\right] \right]^{-1}$$
There is a conductive theorem with the partial condition

▷ Then $\alpha_i \leq \beta_i \leq \gamma_i$ (i.e., α and γ sandwich the partial conditionals)

From Bounds to Simulation

- Simple, exact, and useless method of simulating a draw y from Y
 - \triangleright Draw *r* from $R = R_1, \ldots, R_N$, w/ $R_i \sim U(0, 1)$

▷ For $i \in 1, ..., N$, let $y_i = 1$ if $r_i < \Pr(Y_i = 1 | \theta, Y_{<i}^c = y_{<i}^c) = \beta_i$, else $y_i = 0$

• Problem: we don't know $\beta_i!$ But sometimes, we don't have to...

 \triangleright If $r_i < \alpha_i$, then $r_i < \beta_i$, so we know that $y_i = 1$

 \triangleright If $r_i \geq \gamma_i$, then $r_i \geq \beta_i$, so we know that $y_i = 0$

 $\triangleright~$ In these cases, we say that $lpha,\gamma$ "fix" y_i

Further observation: when y_i is "fixed," this restricts later α, γ values

- $\triangleright \ \alpha \ {\rm and} \ \beta \ {\rm can} \ {\rm only}$ "tighten" as more values are fixed
- \triangleright "Fixation cascade:" fixing y_i tightens bounds for y_{i+k} , fixing it (which may in turn fix others)

Sometimes, will fail; may have to get out and push

► This suggests an algorithm....

The "Bound Sampler"

- Basic algorithm for approximate simulation of $Y \ddagger$
 - \triangleright For $i \in 1, \ldots, N$, perform the following steps:
 - \diamond Initialization: compute α_i, γ_i given $y_{< i}, \theta$
 - \diamond *Fixation*: if $r_i < \alpha_i$, set $y_i = 1$; if $r_i \ge \gamma_i$, set $y_i = 0$
 - ♦ *Perturbation*: if $\alpha_i \leq r_i < \gamma_i$, then "perturb" by drawing $\tau \sim U(\alpha_i, \gamma_i)$ and set $y_i = I(r_i < \tau_i)$ †
 - \triangleright Resulting y vector is an approximate draw from Y
- Bound sampler has some interesting features
 - Draws are independent (unlike MCMC)
 - ▷ Worst-case time complexity fixed at O(Nf(N)), where f(N) is complexity of the initialization step (can be optimal, if f(N) constant)
 - \triangleright Can obtain quality guarantees: if y_i is fixed initially or by an unperturbed cascade, it *must* equal its "true" value!
 - $\diamond~$ "True" means wrt a coupled exact sampler with same R
 - \diamond Quality is a lower bound: unfixed y_i might still be correct, and fixed y_i certainly correct †



























Application to Network Simulation

Bound sampler can easily be used to simulate network processes

- Write network model in ERG form (in terms of adjacency matrix)
- Vectorize" the adjacency matrix via row or column concatenation
- Apply the bound sampler to the resulting vector
- Potential advantages
 - Fixed execution time (and can be very fast)
 - > Unsupervised (e.g., no convergence diagnostics)
 - > Quality guarantees

Potential drawbacks

- Fast implementation can require smart data structures (need cheap initialization)
- If it doesn't work well, you're screwed (but could use smarter perturbation heuristics); can't trade time for quality

Example: 9/11 PATH Radio Communications



Example: Hunter and Handcock Lazega Model (Uniform Perturbations)







Bound Sampler

MCMC

Example: Hunter and Handcock Lazega Model (Pseudo-marginal Perturbations)



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Bound sampler: an interesting direction for ERG simulation

- Fixed execution time (can be worst-case optimal)
- ▷ Ex ante and ex post quality guarantees
- Unsupervised execution (no convergence checking)

Ongoing issues

- ▷ When is it good enough for particular applications?
- Better perturbation tricks (the real key to success!)
- Smart implementation for typical graph statistics (min/max tracking)



Thanks for your attention!

Further Issues: Quality Assessment

- Using β to threshold R yields an exact draw from Y; the sampler approximates this using the (α, γ) bounds on β , plus perturbations (when bounds are insufficient)
- ► Let y^t|r be the unique draw from Y corresponding to realization r of R; the "quality' of draw y|r from the bound sampler is the similarity of y and y^t
- ► Simple measure: $Q(y) = (N D_H(y, y^t))/N$, where D_H is the Hamming distance
 - ▷ Don't know $D_H(y, y^t)$, but can create an upper bound by treating all y_i not fixed by a perturbation-free cascade as incorrect; this give an ex post lower bound on Q
 - ▷ Computing unconditional values for (α_i, γ_i) leads to the ex ante lower bound on $\mathbf{E}Q(Y)$, $1/N \sum_{i=1}^{N} (\gamma_i - \alpha_i)$
- Can also answer more nuanced questions about y^t using y
 - > Considering all y' consistent with the fixed values of y allows exact bounds on all properties of y^t (but perhaps loose ones!)
 - Given multiple draws, can bound expectations for properties of Y by expectations on upper/lower bounds from bound sampler draws
- Note: Q(y) is not independent of y^t! Resampling until one gets a high-quality draw will introduce bias! (No free lunch, alas....) [Return]

Further Issues: Perturbation

- When we "perturb" the algorithm, we are really estimating the unknown β_i
 - $\triangleright\,$ The threshold, $\tau,$ is our "estimator"
- ► By default, select uniformly between bounds, but can do better
 - ▷ Clearly, $\mathbf{E}\beta = \sum_{i=1}^{N} Y_i / N$, so want to favor low/high τ when Y is sparse/dense
 - ▷ Taking an initial approximation to the pdf of β as a prior, can update based on (α, γ) and draw τ from the posterior
 - Diffuse beta distribution centered on expected density seems to be an easy improvement on the uniform
 - ▷ In theory, could use better approximations that take $y_{<i}$ into account; as $\tau \to \beta_i$, y|r approaches its exact value
 - ♦ Have had good luck w/"pseudo-marginalization" method to approximate β_i by averaging full conditionals over a baseline Bernoulli model
- ► This is the principal route for improving simulation quality clever ideas are welcome!

[Return]



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1: for i in 1, ..., N do
         Set \alpha'_i := \min_{y' \in \mathcal{Y}_N : y'_{< i} = y_{< i}} \Pr\left(Y_i = 1 \left| \theta, Y_i^c = y'_i^c\right.\right)
  2:
         Set \gamma'_i := \max_{y' \in \mathcal{Y}_N : y'_{< i} = y_{< i}} \Pr\left(Y_i = 1 \left| \theta, Y_i^c = y'_i^c\right.\right)
  3:
         Draw r_i from U(0,1)
  4:
         if r_i < \alpha'_i then
  5:
            Set y_i = 1
  6:
         else if r_i \geq \gamma'_i then
  7:
             Set y_i = 0
  8:
         else
  9:
             Draw \tau from U(\alpha'_i, \beta'_i)
10:
11:
            if r_i < \tau then
            Set y_i := 1
12:
            else
13:
                Set y_i := 0
14:
             end if
15:
         end if
16:
17: end for
18: return y
[Return]
```

1 References

Butts, C. T. (2010). Bernoulli graph bounds for general random graphs. Technical Report MBS 10-07, Irvine, CA.